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A Scalable Multicast Scheme for Distributed DAG Scheduling

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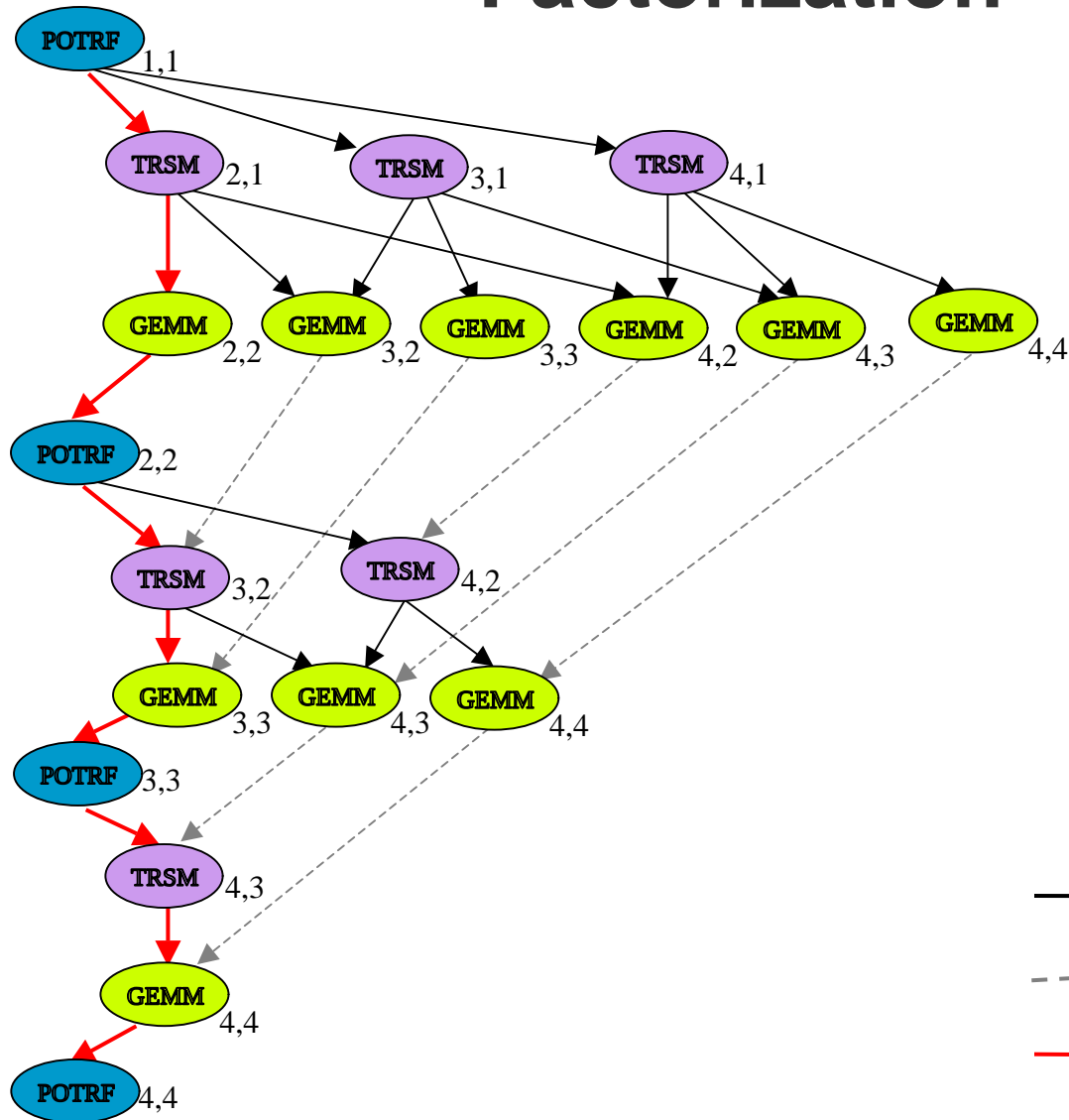
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Motivation

- High performance on multicore machines
- New software should have two characteristics:
 - Fine grain threads
 - Asynchronous execution
- We want to use dynamic DAG scheduling
- Extremely Scalable
 - We are thinking of millions of processing cores.
 - Distributed-memory

A DAG Example for Cholesky Factorization



1,1			
2,1	2,2		
3,1	3,2	3,3	
4,1	4,2	4,3	4,4

- loop indep. dependency
- - - → loop carried dependency
- critical path (T_∞)

Simple Programming Model

- Symbolic DAG interface:
 - `int get_num_parents(const Task t);`
 - `int get_children(const Task t, Task children);`
 - `set_entry_task(const Task t);`
 - `set_exit_task(const Task t);`

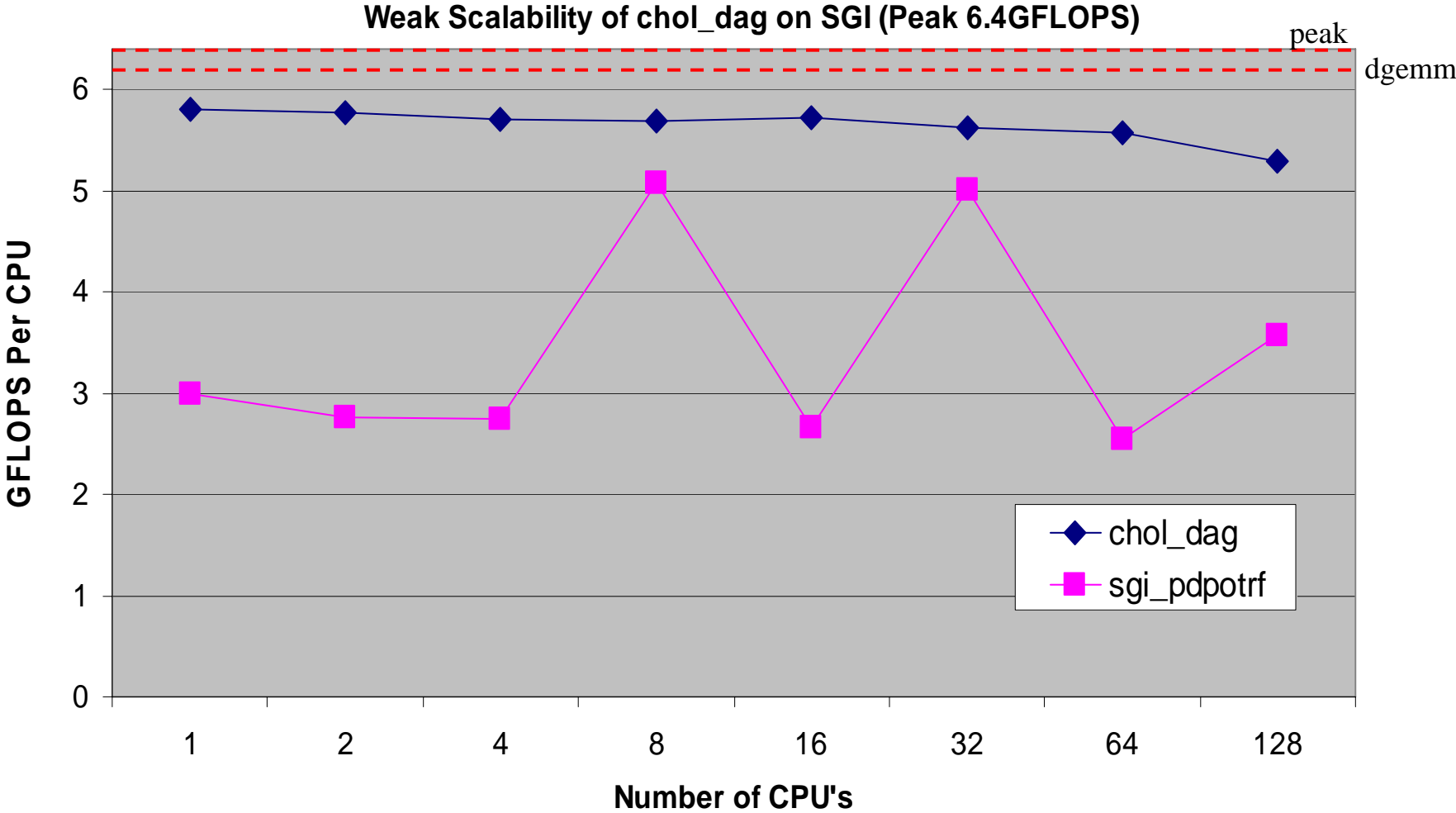
Interface Definition for Cholesky Factorization

```
struct Task {  
    int type; // what task  
    int k;    // iteration index  
    int i, j; // row, column index  
    int priority;  
}
```

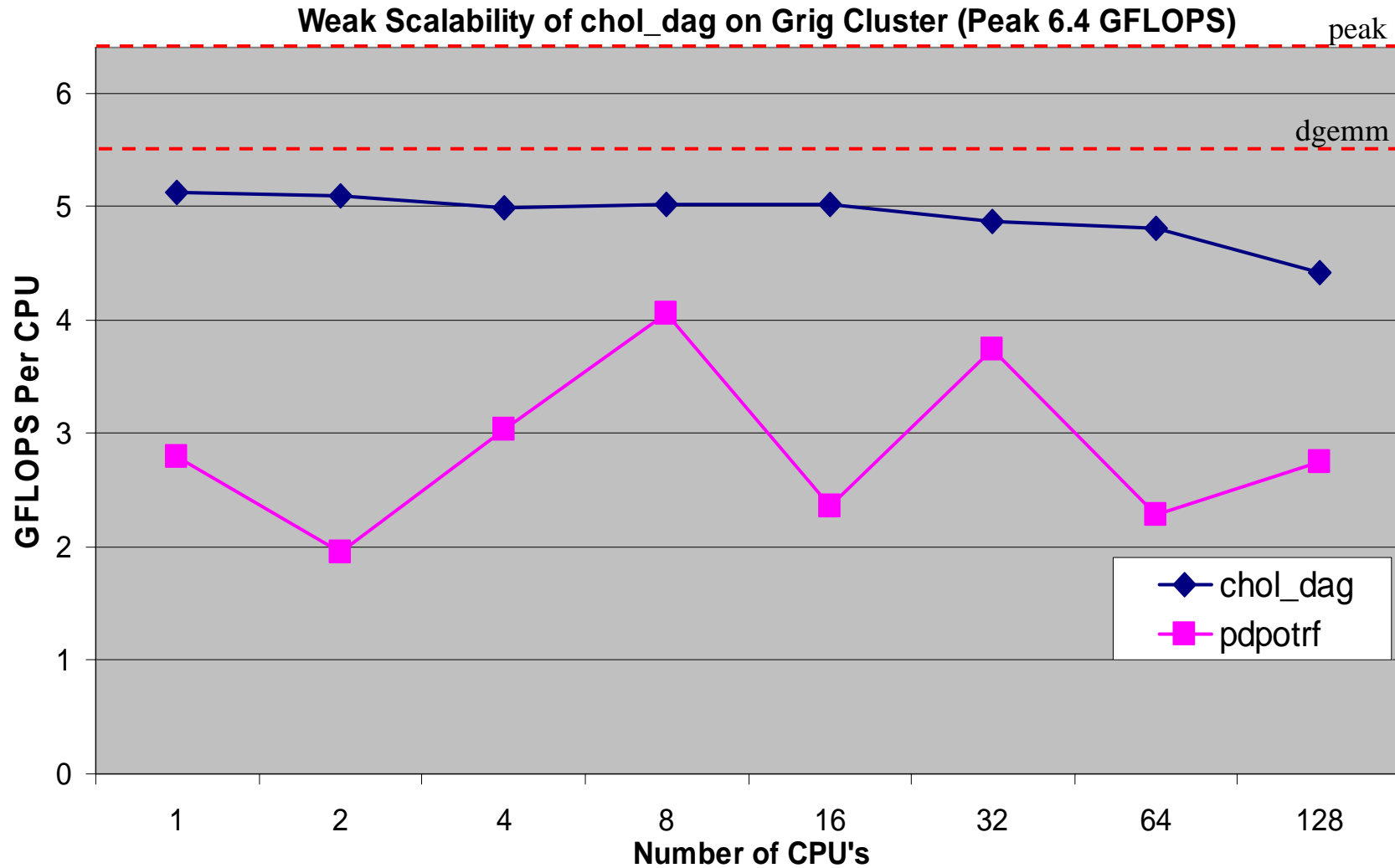
```
int get_num_parents(Task t) {  
    if (t.type = POTRF) return 1  
    if (t.type = TRSM) return 2  
    if (t.type = GEMM) {  
        if (diagonal) return 2  
        else return 3  
    }  
}
```

```
int get_children(Task p, Task* buf, int nblks) {  
    if (p.type = POTRF) {  
        /* along p's column but below p */  
        buf := {TRSM task t | t.j = p.j & t.i ∈ ( p.i, nblks]}  
    }  
    if (p.type = TRSM) {  
        /* a row and a column (both with index p.i) */  
        buf := {GEMM task t | t.i = p.i & t.j ∈ (p.j, p.i] or  
                t.j = p.i & t.i ∈ [p.i, nblks] }  
    }  
    if (p.type = GEMM) {  
        /* has a single child */  
        if (diagonal) buf := a POTRF task  
        else if (below diag) buf := a TRSM task  
        else buf := a GEMM task  
    }  
    return |buf|  
}
```

Performance on SGI Altix 3700 BX2



Performance on the Grig Cluster





Multicast Scheme Overview

- Application-level routing layer
- Hierarchical abstraction of a system
- Each process has a topology ID.
 - Like zip code
 - The longer the common prefix of two topo_ids, the closer they are.
- Compact routing table
 - An extension to Plaxton's neighbor table [1]

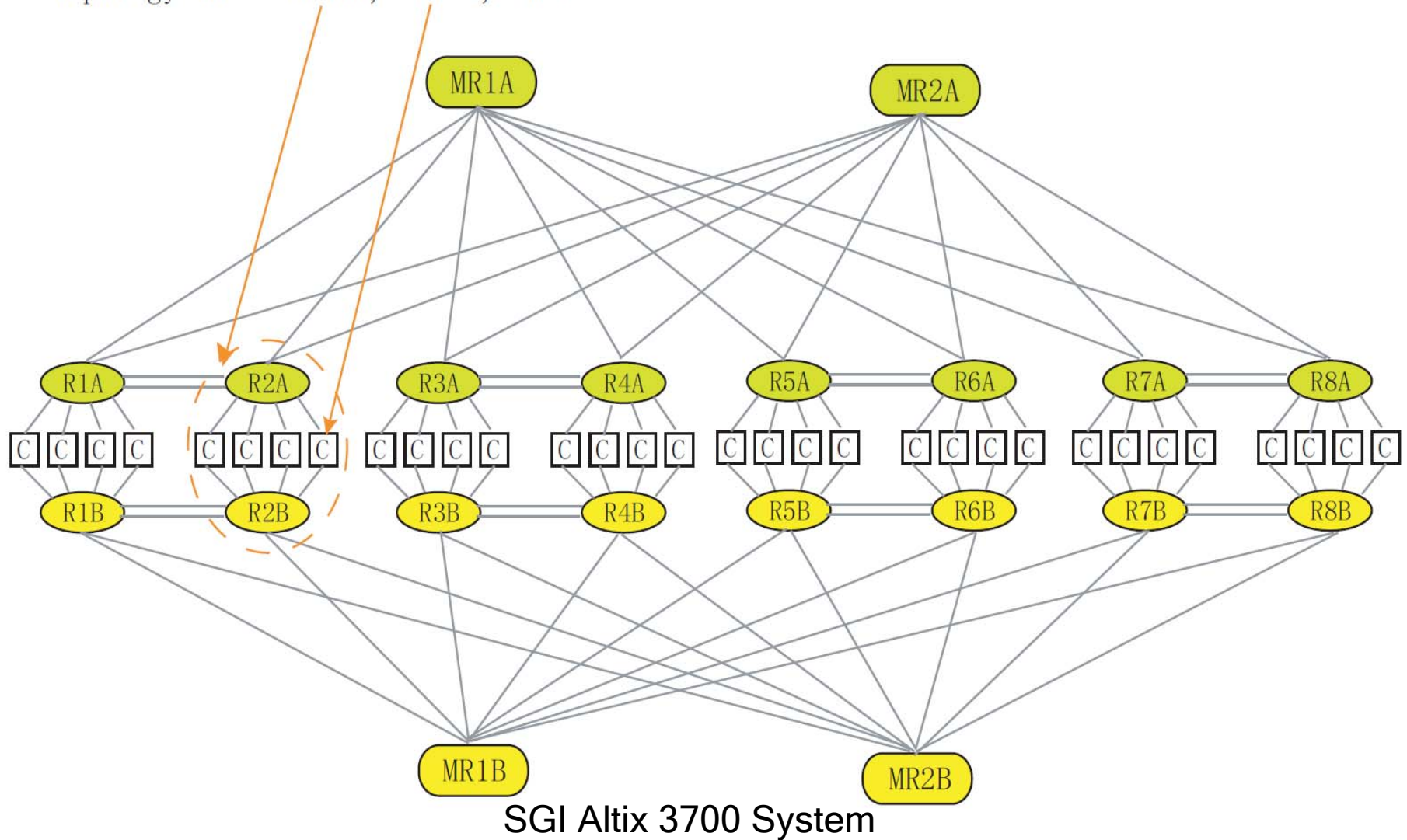
[1] Plaxton, C. G., Rajaraman, R., and Richa, A. W. 1997. Accessing nearby copies of replicated objects in a distributed environment. SPAA '97.

Topology ID

- Assign IDs to the whole system (i.e., T_{system})
 - T_{program} of a user program $\subset T_{\text{system}}$
- A topology ID is a number of digits.
 - E.g., 256 nodes consist of 4 digits with base 4.

 - E.g., 2048 nodes consist of 4 digits with base 8.

- We assume that two nodes with a longer common prefix are closer on the physical network.

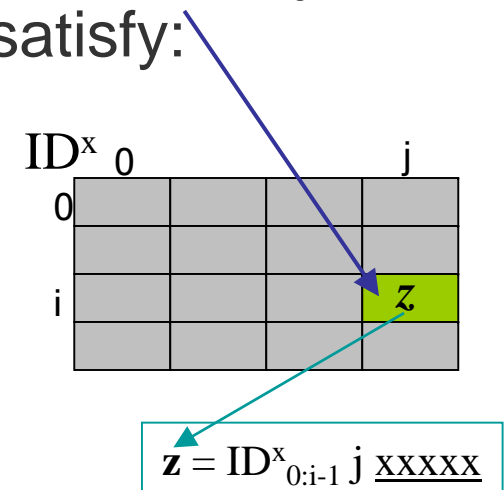
Topology ID Example

Topology ID = <3bits, 2bits, 2bits>



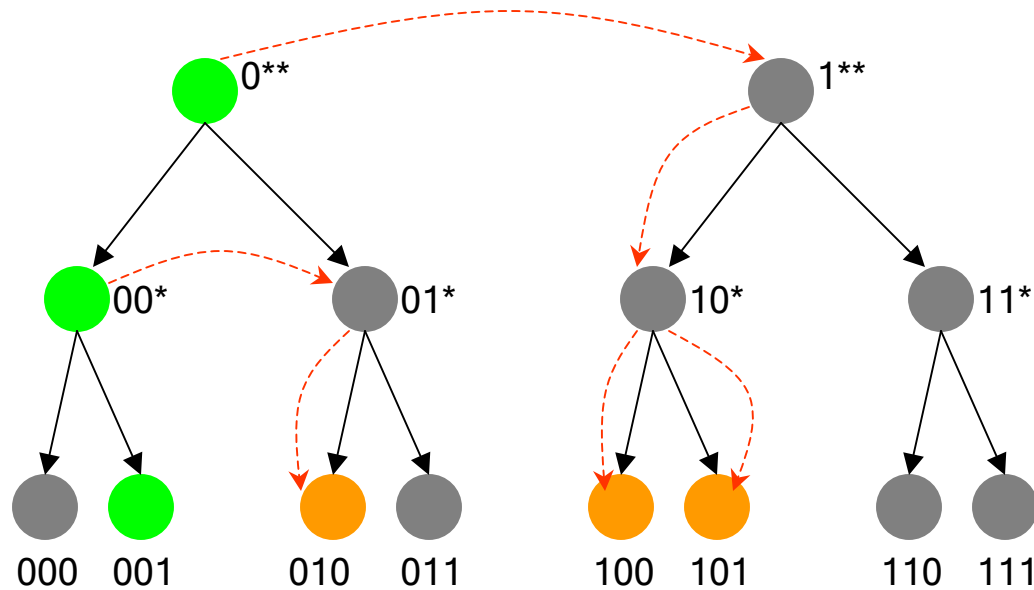
Compact Routing Table

- Suppose process x has a routing table and $\text{Table}[i,j]$ stores process z , then ID^x and ID^z must satisfy:
 - $\text{ID}^x_0 \text{ID}^x_1 \dots \text{ID}^x_{i-1} = \text{ID}^z_0 \text{ID}^z_1 \dots \text{ID}^z_{i-1}$,
 - $\text{ID}^z_i = j$ (i.e., $(i+1)$ th digit of $\text{ID}^z = j$).
- Routing table could have empty entries.
- Always search for the forwarding host on the $\text{LCP}(x,y)$ row.
- At most $(\log_2 P)/(\text{base})$ steps
- $O(\lg P)$ space cost
 - 1 million cores \rightarrow 80(5x16) entries
 - 1 billion cores \rightarrow 192(6x32) entries



A Multicast Example

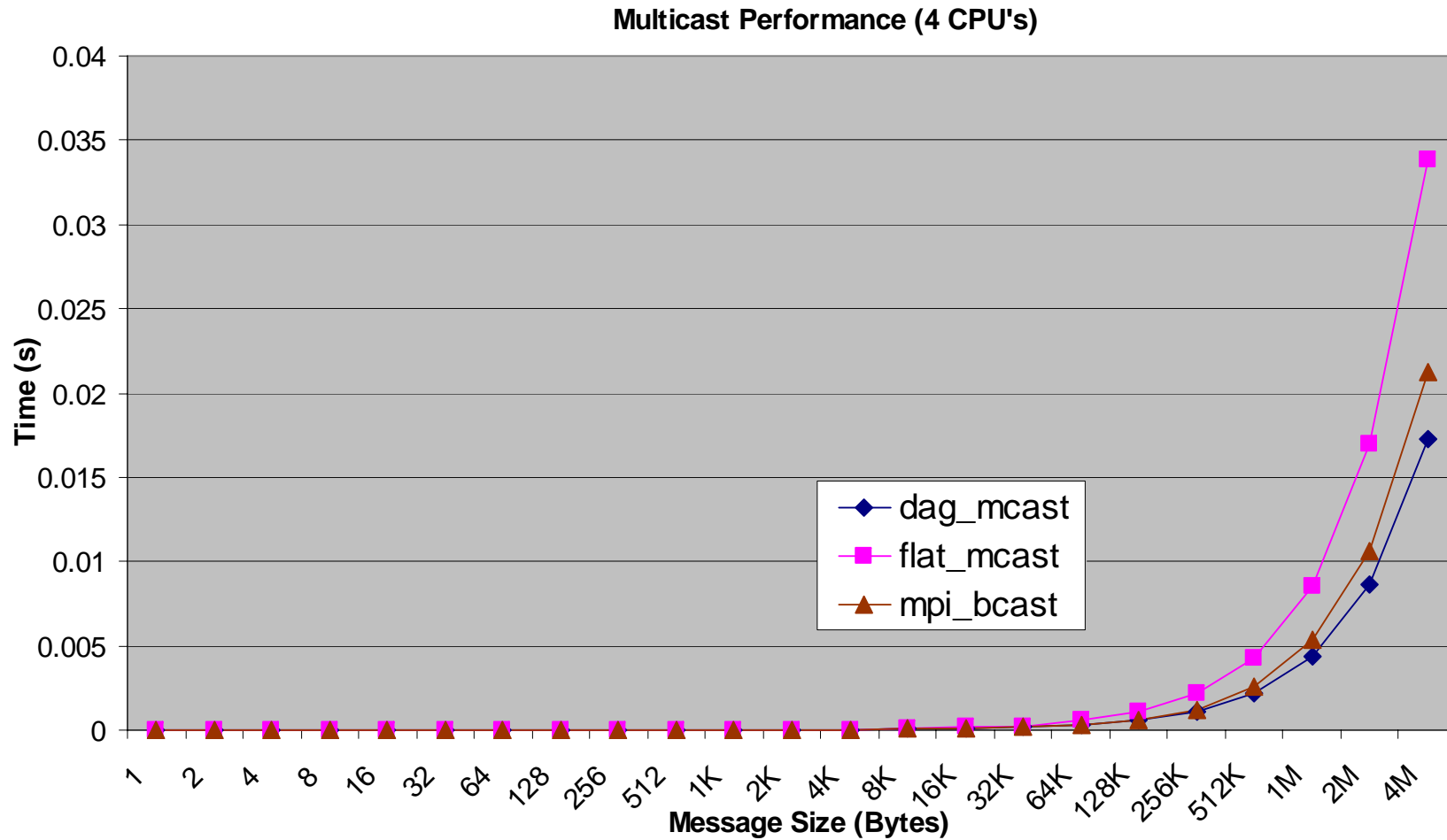
- Node 001 multicasts data to nodes {010, 100, 101}.



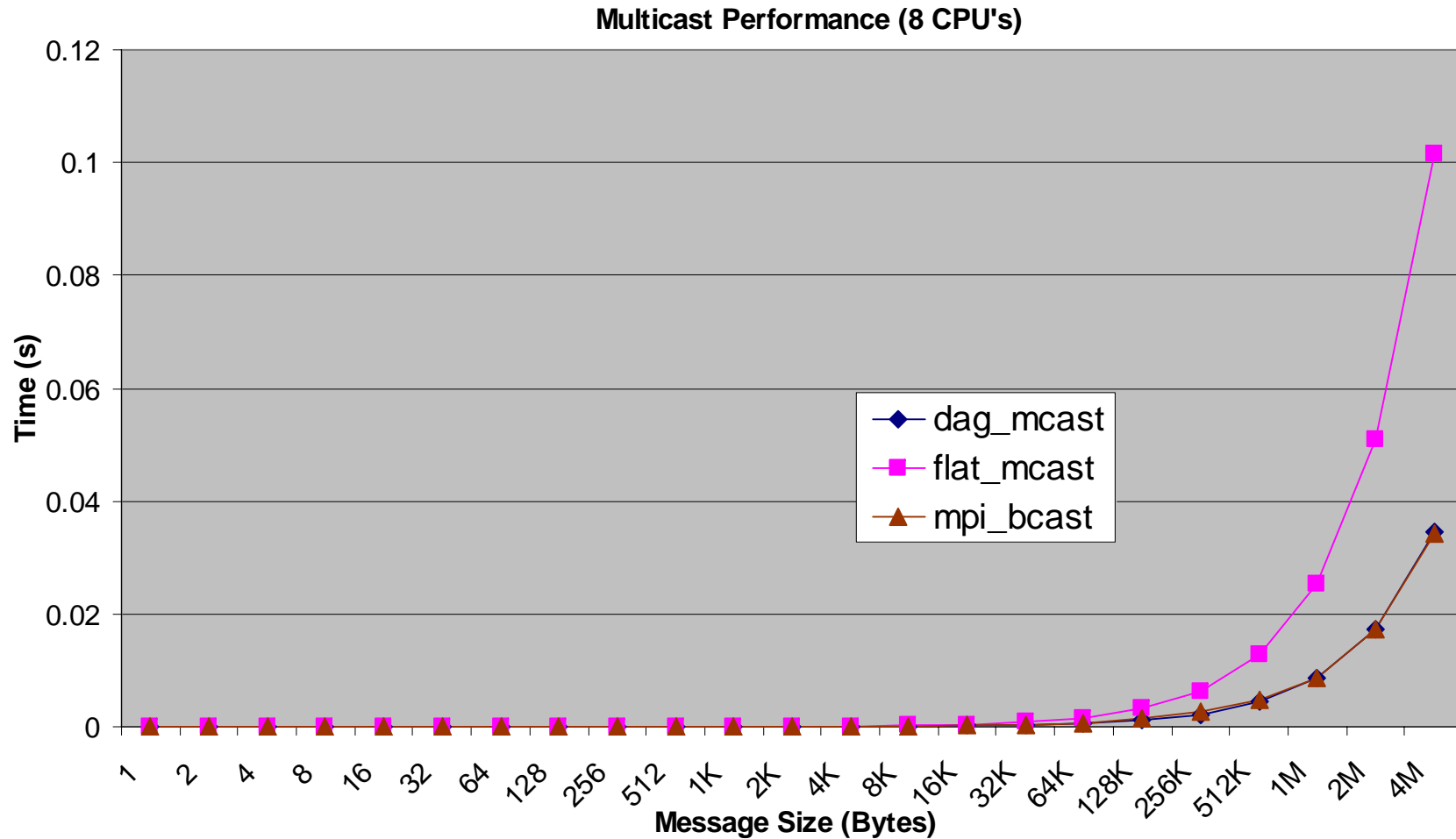
Grig Cluster

- grig.sinrg.cs.utk.edu
- 64 nodes, dual-CPU per node
 - Intel Xeon 3.20GHz
 - Peak performance 6.4 GFLOPS
 - **Myrinet interconnection** (MX 1.0.0)
- Goto BLAS 1.26
 - DGEMM performance 5.57 GFLOPS
 - 87% of peak performance (upper bound)
- **MPICH-MX 1.x**
- gcc 64 bits

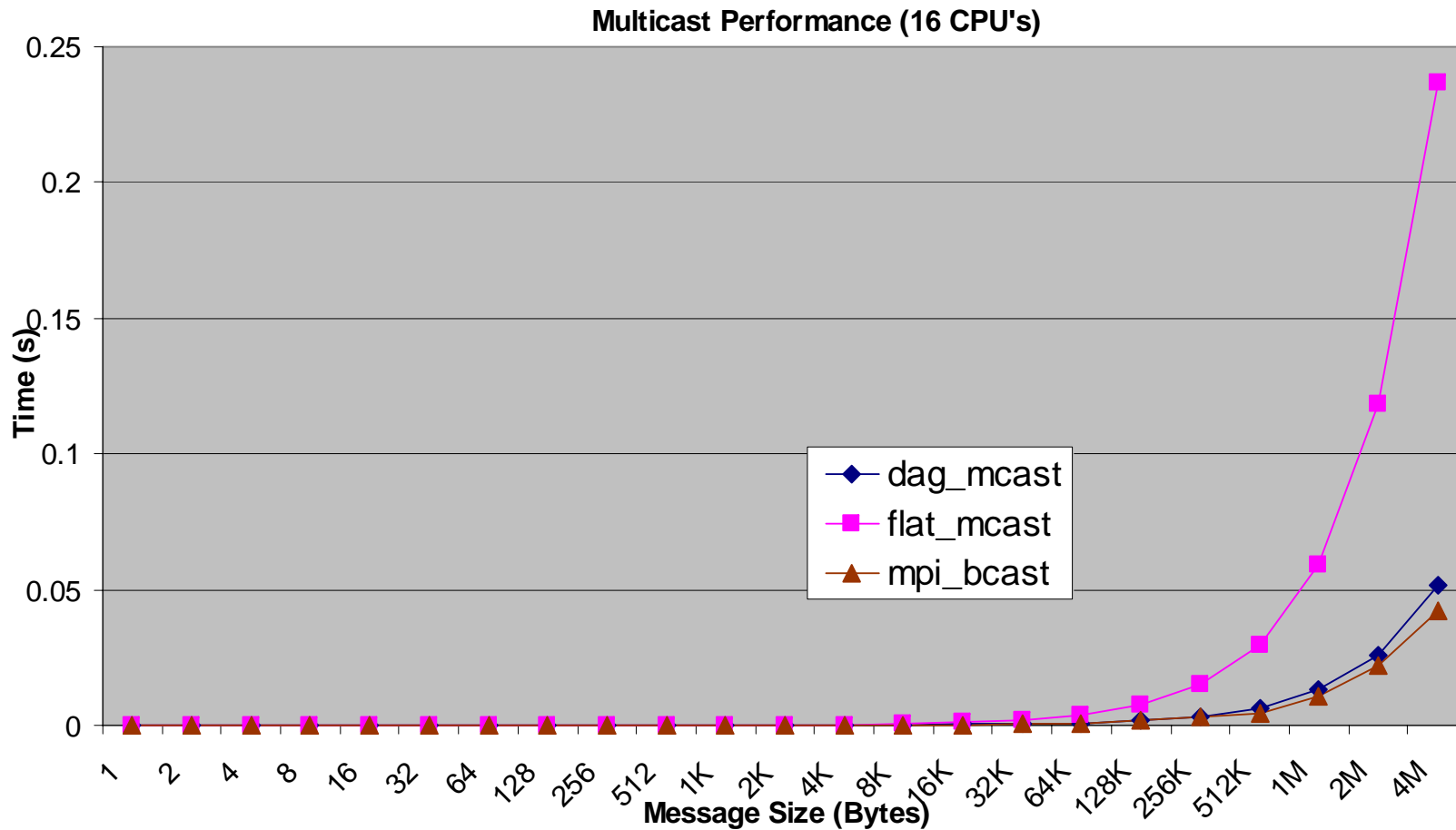
Multicast on a Cluster (4 CPUs)



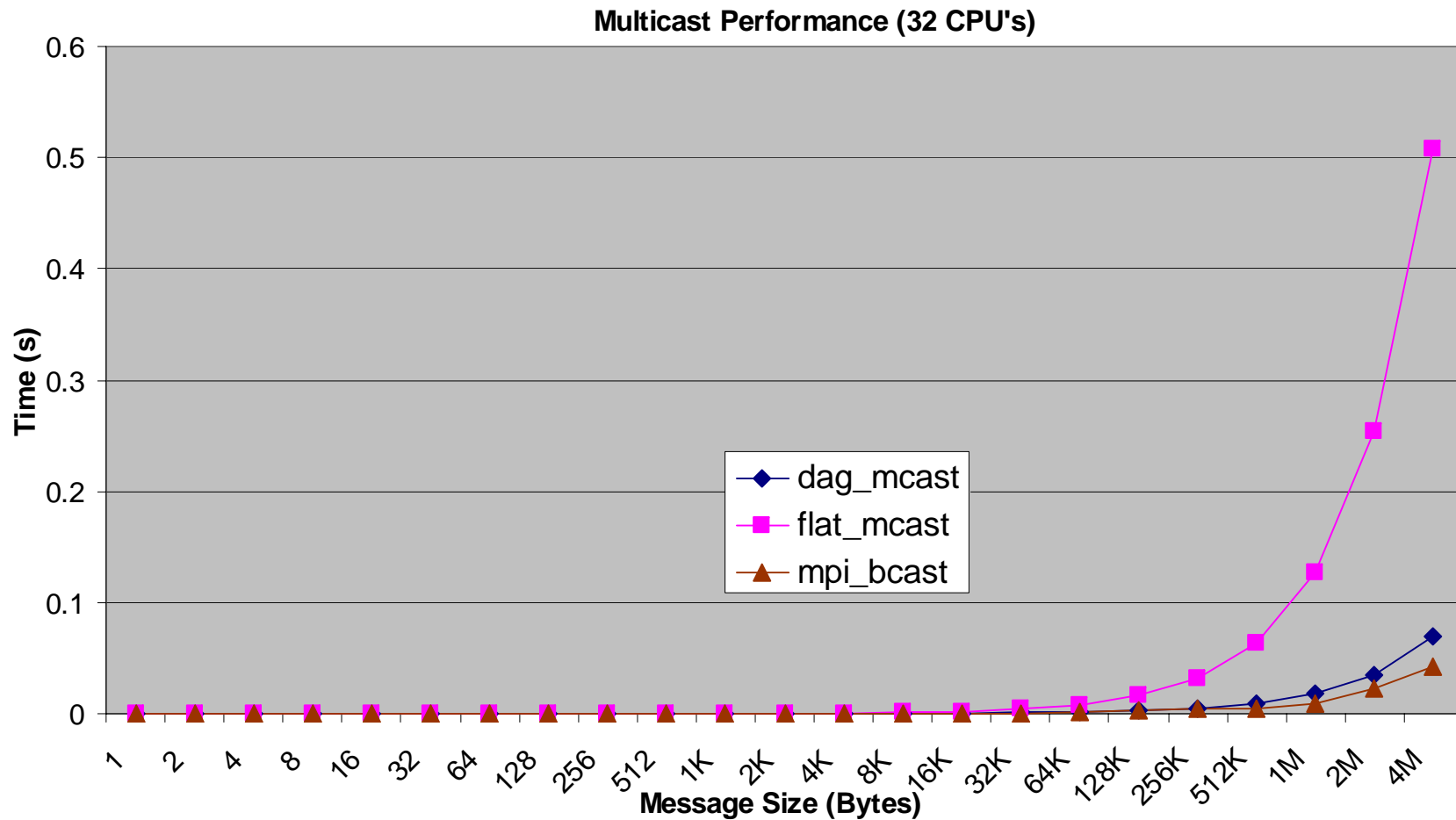
Multicast on a Cluster (8 CPUs)



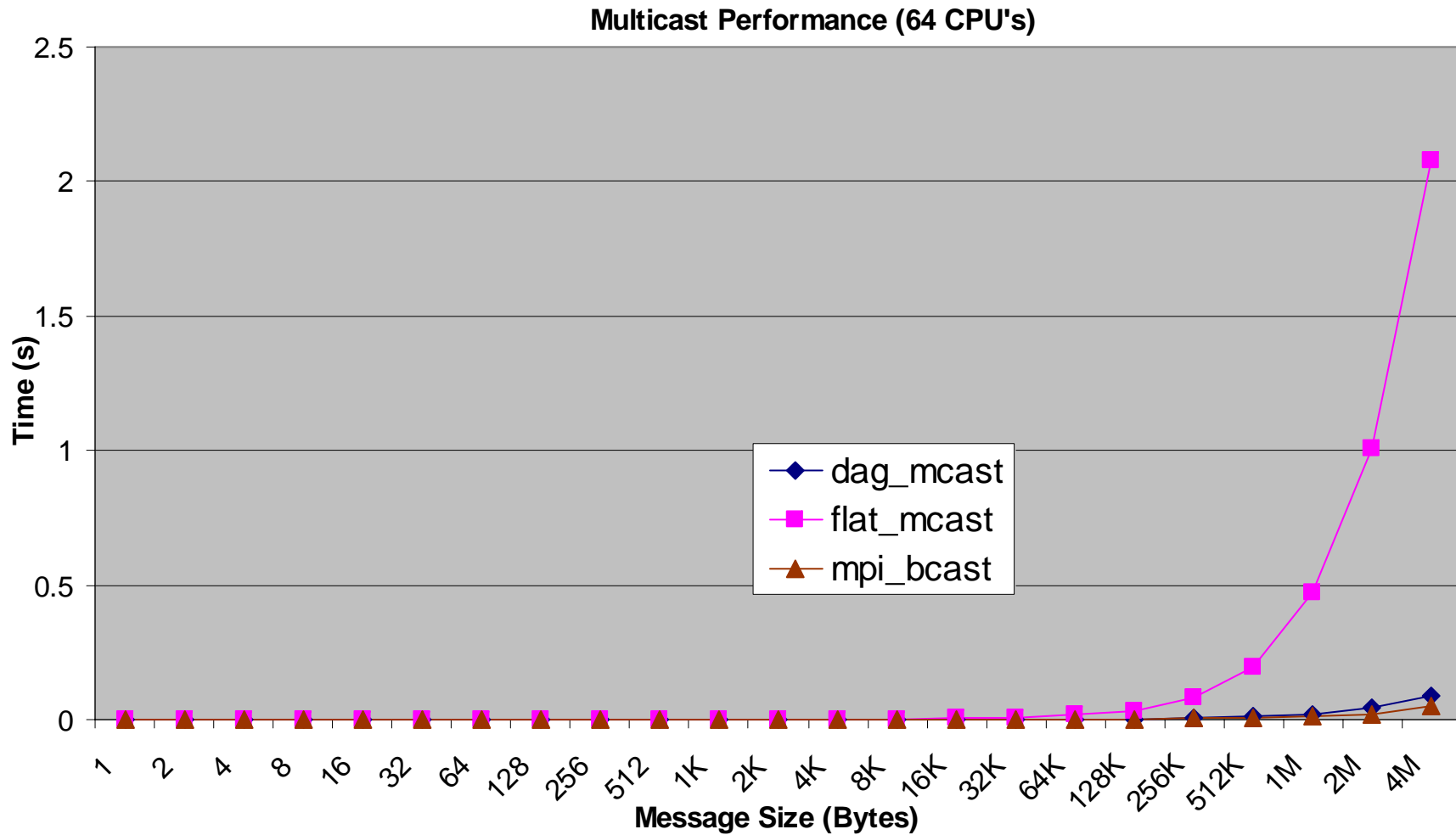
Multicast on a Cluster (16 CPUs)



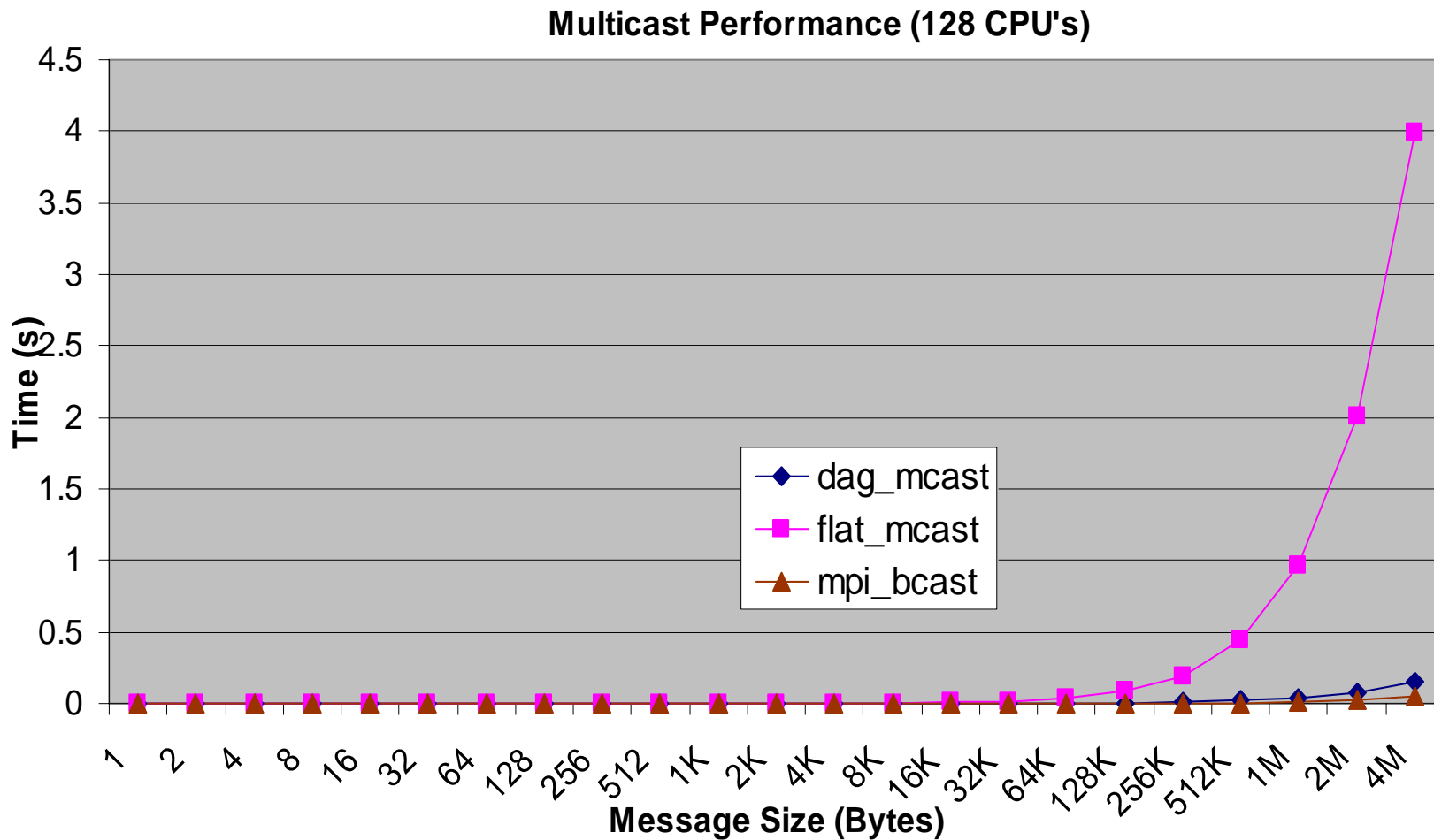
Multicast on a Cluster (32 CPUs)



Multicast on a Cluster (64 CPUs)



Multicast on a Cluster (128 CPUs)



Summary

- Support scalable multicast in distributed DAG scheduling
- Important features:
 - Non-blocking
 - Topology-aware
 - Scalable in terms of routing-table space and #steps
 - Dead-lock free
 - No requirement of communication group creation
 - Support multiple concurrent multicasts
- Performance is close to vendor's MPI_Bcast.