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# Toward autonomic QoS in Grid-aware applications: the ASSIST experiment



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## Outline

### • Motivating ...

- high-level programming for the grid
- application adaptivity for the grid
- ASSIST basics & adaptivity in ASSIST
  - mechanisms
  - demo & some experiments
- Components & QoS
  - autonomic managers
  - QoS contracts
- Concluding remarks



## // progr. & the grid

- concurrency exploitation, concurrent activities set up, mapping/scheduling, communication/synchronization handling and data allocation, ...
- manage resources heterogeneity and unreliability, networks latency and bandwidth unsteadiness, resources topology and availability changes, firewalls, private networks, reservation and jobs schedulers, ...

... and a non trivial QoS for applications
not easy leveraging only on middleware
D. Gannon et al. opened the way (GrADS@Rice)



## ASSIST idea

"moving most of the Grid specific efforts needed while developing highperformance Grid applications from programmers to grid tools and run-time systems"



ASSIST is a high-level programming environment for grid-aware // applications. Developed at Uni. Pisa within several national & EU projects. First version in 2001. Open source under GPL.

## app = graph of modules

Programmable, possibly nondeterministic input behaviour

P2

Sequential or parallel module

P1

input

Typed streams of data items

P3

P4



output

### native + standards













An "input section" can be programmed in a CSP-like way

Data items can be distributed (scattered, broadcasted, multicasted) to a set of

#### **Virtual Processes**

which are named accordingly to a topology





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**V**P

VP

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while(...)
 forall VP(in, out)
 barrier

data is logically shared by VPs (owner-computes)



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Easy to express standard paradigms (skeltons), such as farm, deal, haloswap, map, apply-to-all, forall, ...

## parmod implementation





## Compiling & running



## Application adaptivity

- Adaptivity aims to dynamically control program configuration (e.g. parallel degree) and mapping
  - for performance (high-performance is a natural subtarget)
  - for fault-tolerance (enable to cope with unsteadiness of resources, and some kind of faults)



## Adaptivity recipe (ingredients)

### 1. Mechanism for adaptivity

- reconf-safe points
  - in which points a parallel code can be safely reconfigured?
- reconf-safe point consensus
  - different parallel activities may not proceed in lock-step fashion
- add/remove/migrate computation & data
- 2. Managing adaptivity
  - QoS contracts
    - Describing high-level QoS requirement for modules/applications
  - "self-optimizing" modules/components
    - under the control of an autonomic manager



## Mechanisms

#### • At parmod level

- add/remove/migrate VPs
- very low-overhead due to knowledge coming from high-level semantics + suitable compiling tools
- At component level
  - create/destroy/wire/unwire parallel entities
  - medium/large overhead due to underlying API for staging, run, ...
- Not addressed in this talk (see references in the paper: Europar 05, ParCo 05, ...), I just show a short demo













## overhead? (mSecs)

parmod kind	Data-parallel (with shared state)					Farm (without shared state)						
reconf. kind	add PEs			remove PEs			add PEs			remove PEs		
# of PEs involved	1→2	2→4	4→8	2→1	4→2	8→4	1→2	2→4	4→8	$2 \rightarrow 1$	4→2	8→4
$egin{array}{c} R_l & { m on-barrier} \ R_l & { m on-stream-item} \end{array}$	1.2 4.7	1.6 12.0	2.3 33.9	0.8 3.9	1.4 6.5	3.7 19.1	$\sim 0$	$\sim 0$	$\sim 0$	$\sim 0$	$\sim 0$	$\sim \overline{0}$
$R_t$	24.4	30.5	36.6	21.2	35.3	43.5	24.0	32.7	48.6	17.1	21.6	31.9

GrADS papers reports overhead in the order of hundreds of seconds (K. Kennedy et al. 2004), this is mainly due to the stop/restart behavior, not to the different running env.



## Autonomic Computing



- AC emblematic of a vast hierarchy of selfgoverning systems, many of which consist of many interacting, self-governing components that in turn comprise a number of interacting, self-governing components at the next level down.
- IBM "invented" it in 2001 (control with selfawareness, from human body autonomic nervous system)
  - self-optimization, self-healing, selfprotection, self-configuration = selfmanagement
- control loop, of course, exists from mid of last century



## Autonomic behavior



- monitor: collect execution stats: machine load, VPM service time, input/output queues lenghts, ...
- analyze: instanciate performance models with monitored data, detect broken contract, in and in the case try to indivituate the problem
- plan: select a (predefined or user defined) strategy to reconvey the contract to valid status. The strategy is actually a list of mechanism to apply.
- execute: leverage on mechanism to apply the plan



## Autonomic behavior

#### Managed element (module, component)

Monitor

QoS data

Plan

Analyze

ecute

broken

contract

next

configuration

 monitor: collect queues lenghts,

- analyze: instanci contract, in and i
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Autonomic behavior as been included in NGG2/3 (Next Generation Grid) EU founding recommendation as prerequisite for Grid computing

service time, input/output

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- ASSIST graphs can be enclosed in components
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- they can be wired to other legacy components (e.g. CCM)
- currently *native component model*, already converging in the forthcoming GCM (authors involved in CoreGRID NoE, WP3)



## managed components



- modules and components are controlled by managers
- managers implements NF-ports



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- modules and components are controlled by managers
- managers implements NF-ports
- the distributed coordination of managers enable the managing of the application as whole (the top manager being the Application Manager)



### QoS contract (of the experiment I'll show you in a minute)

Perf. features	$QL_i$ (input queue level), $QL_o$ (input queue level), $T_{ISM}$ (ISM service time), $T_{OSM}$ (OSM service time), $N_w$ (number of VPMs), $T_w[i]$ (VPM <sub>i</sub> avg. service time), $T_p$ (parmod avg. service time)
Perf. model	$T_p = \max\{T_{ISM}, \sum_{i=1}^n T_w[i]/n, T_{OSM}\}, T_p < K \text{ (goal)}$
Deployment	$\operatorname{arch} = (i686\text{-pc-linux-gnu} \lor \operatorname{powerpc-apple-}$

Adapt. policy goal\_based

darwin\*)

## experiment: stateless farm



- contract:
  - keep a given service time
  - contract change along the run

## Experimenting heterogeneity





## Experimenting heterogeneity



Not only Intel+linux: similar experiments has been run on Linux, Mac, Win, and a mixture of them



#### Data-par experiment (STP) 25% 35% 36% 31% 32% 41% 51% 25% 22% 23% **Distribution of load among platforms (n. of VPs)** 400 D 300 С 200 B 100 A 0 80% 40% **Relative Unbalance** 0% 6 **Iteration time** 3 50 150 200 250 300 350 100 400 Time (iteration no.) 67

## Conclusions 1/2

- Application adaptivity in ASSIST
  - complex, but trasparent (no burden for the programmers)
    - they should just define they QoS requirements
    - QoS models are automatically generated from program structure (and don't depend on seq. funct.)
  - dynamically controlled, efficiently managed
    - catch both platforms unsteadiness and code irregular behavior in running time
    - performance models not critical, reconfiguration does not stop the application
    - key feature for the grid



## Conclusions 2/2

### ASSIST cope with

- grid platform unsteadiness
- interoperability with standards
  - and rely on them for many features
- high-performance
- app deployment problems on grid
  - private networks, job schedulers, firewalls, ...
- QoS of the whole application through hierarchy of managers





ASSIST is open source under GPL http://www.di.unipi.it/Assist.html